

# Effective Parallelization of the Vehicle Routing Problem

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# Capacitated Vehicle Routing Problem (CVRP)

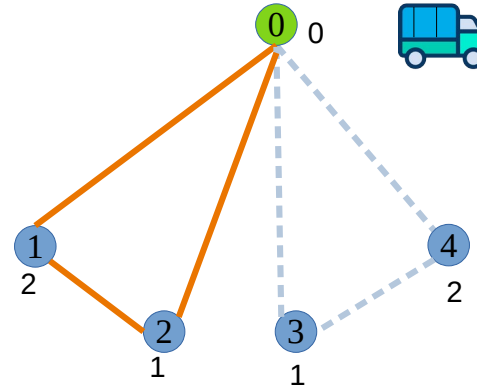
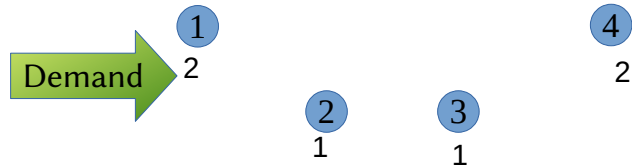
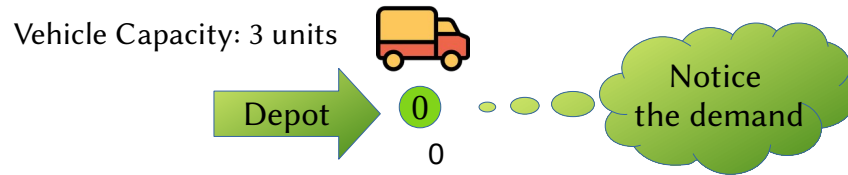




**Input** : Given  $n$  nodes (single Depot and customers) with their coordinates  $(x_i, y_i)$  and demands  $d_i > 0$  for  $i \in n$ , Vehicle capacity  $C$ . Node 0 is Depot and has zero demand.

**Output**: Set of routes serving all the customers respecting the vehicle capacity from/to Depot.

**Goal** : Minimize total distance travelled.

If Capacity = sum of demands  $d_i$ ,  
CVRP  $\rightarrow$  Travelling Salesman Problem



	route1	$\begin{matrix} 0 & 1 & 2 & 0 \end{matrix}$
	route2	$\begin{matrix} 0 & 3 & 4 & 0 \end{matrix}$



# Motivation

## Current state-of-the-art

- work only on smaller instances
- has a large solution Gap.
- takes a lot of time.

Instance	Number of customers	Time (s)	
		Base2	Base1
Flanders2	30,000	8,355	2,534
Flanders1	20,000	7,768	2,031
Brussels1	15,000	7,164	871

**Table 1: State-of-the-art GPU methods are time-consuming.**



**RQ1.** Can we invent a simpler algorithm?



**RQ2.** Can we reduce Gap on large instances?



**RQ3.** Design Parallelization friendly algorithms?

## Our ParMDS

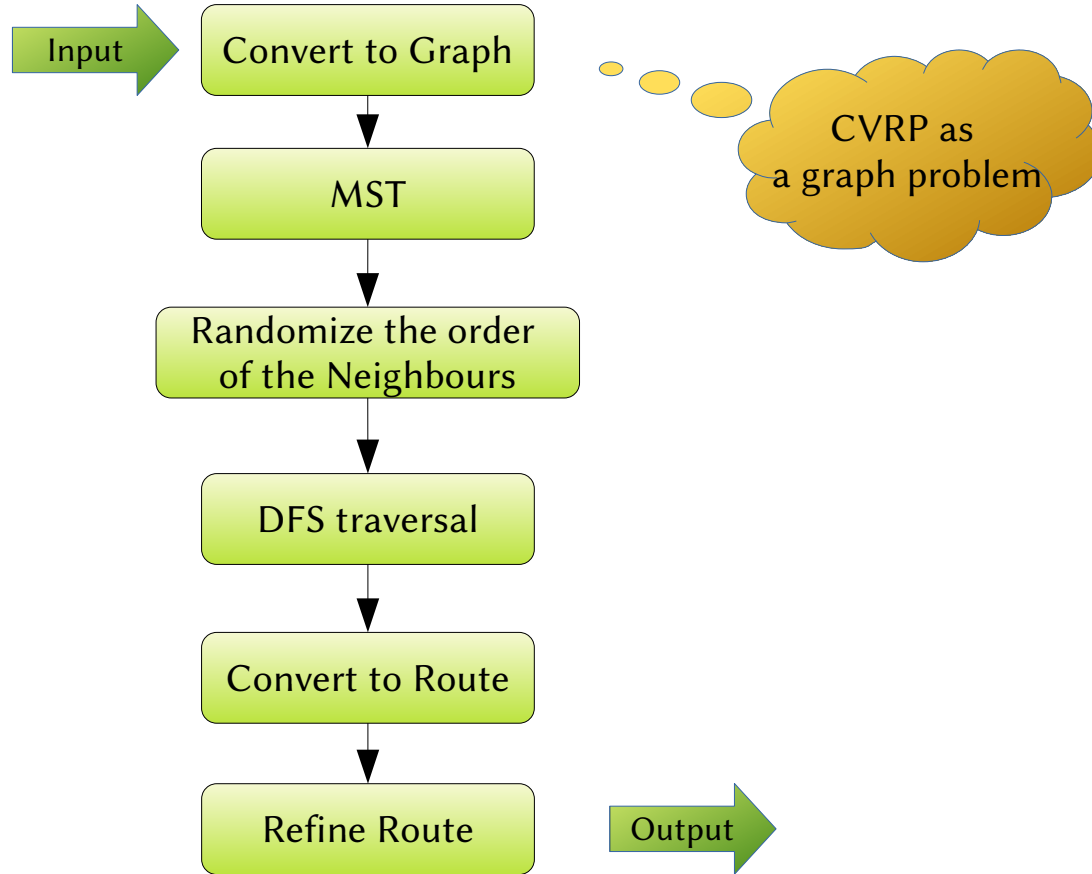
- Serial and **Par**allel implementation
- Combining **M**ST and **D**FS
- Uses Local-search approach
- Uses Randomization approach

$$\text{Gap} = \frac{Z_S - Z_{BKS}}{Z_{BKS}} \times 100$$

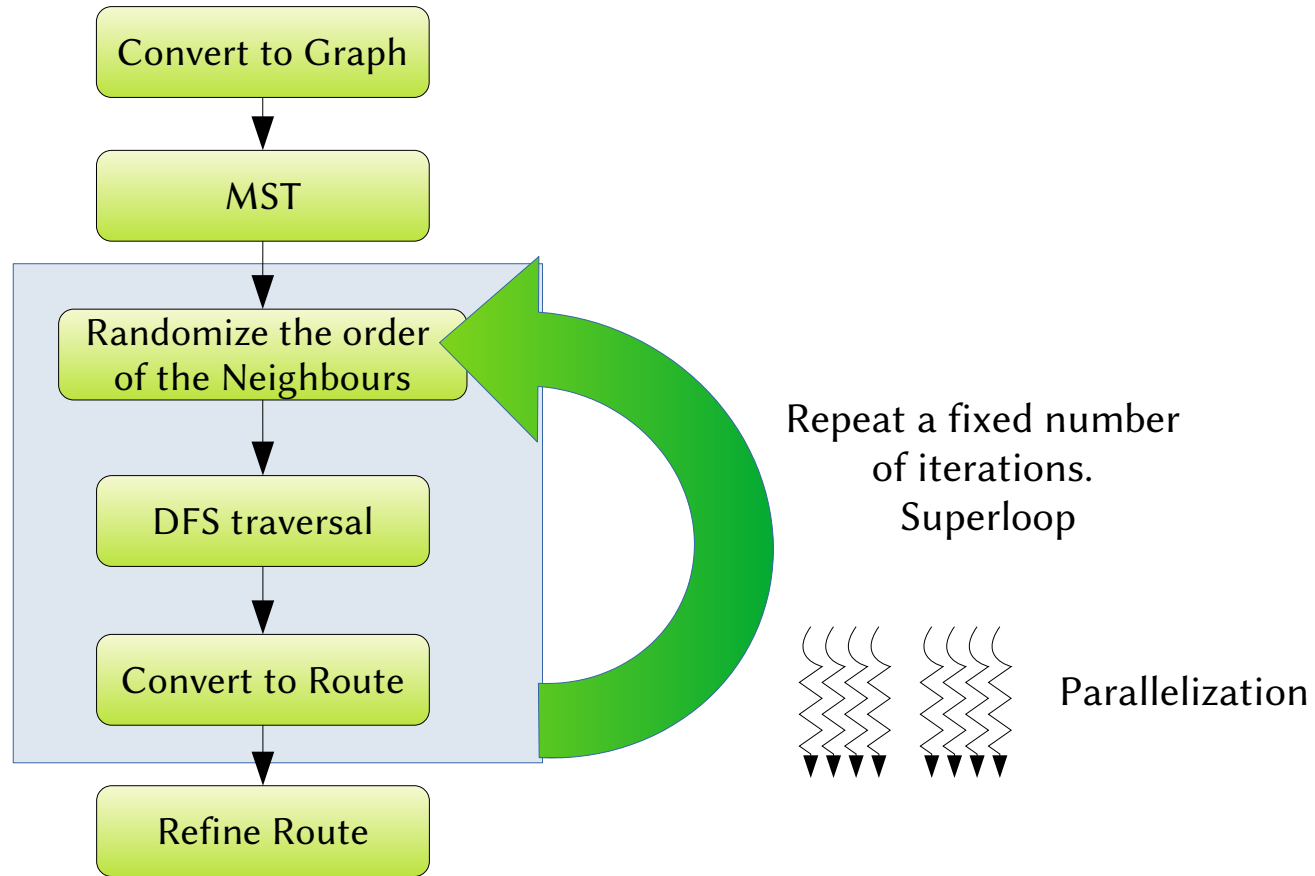
**Baseline1:** P. Yelmewad and B. Talawar. Parallel Version of Local Search Heuristic Algorithm to Solve Capacitated Vehicle Routing Problem, Cluster Computing, 2021.

**Baseline2:** M. Abdelatti and M. Sodhi. An improved GPU-accelerated heuristic technique applied to the Capacitated Vehicle Routing Problem, GECCO, 2020.

# Overview - ParMDS



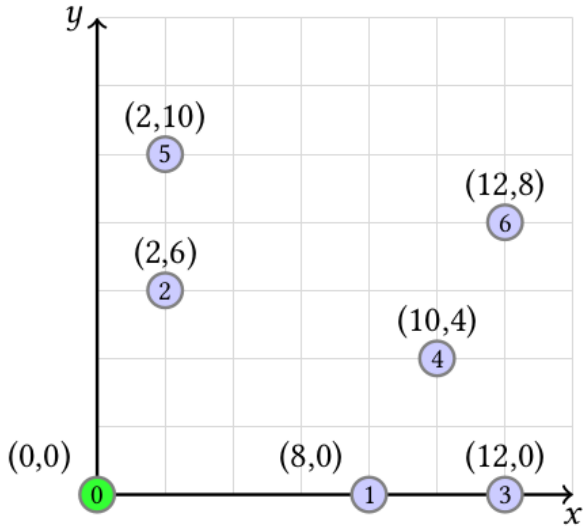
# Overview - ParMDS



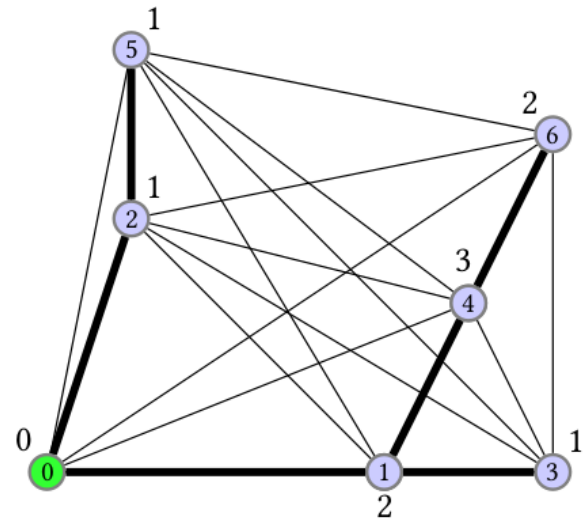
# Example - Overview



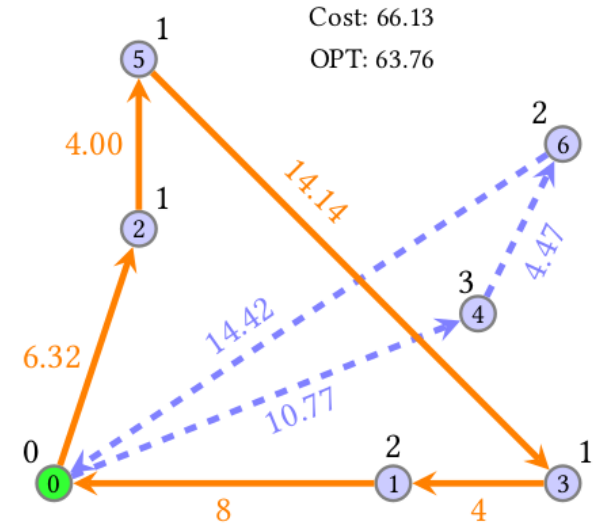
Closer to Optimal Cost



(a) Input instance  $I$



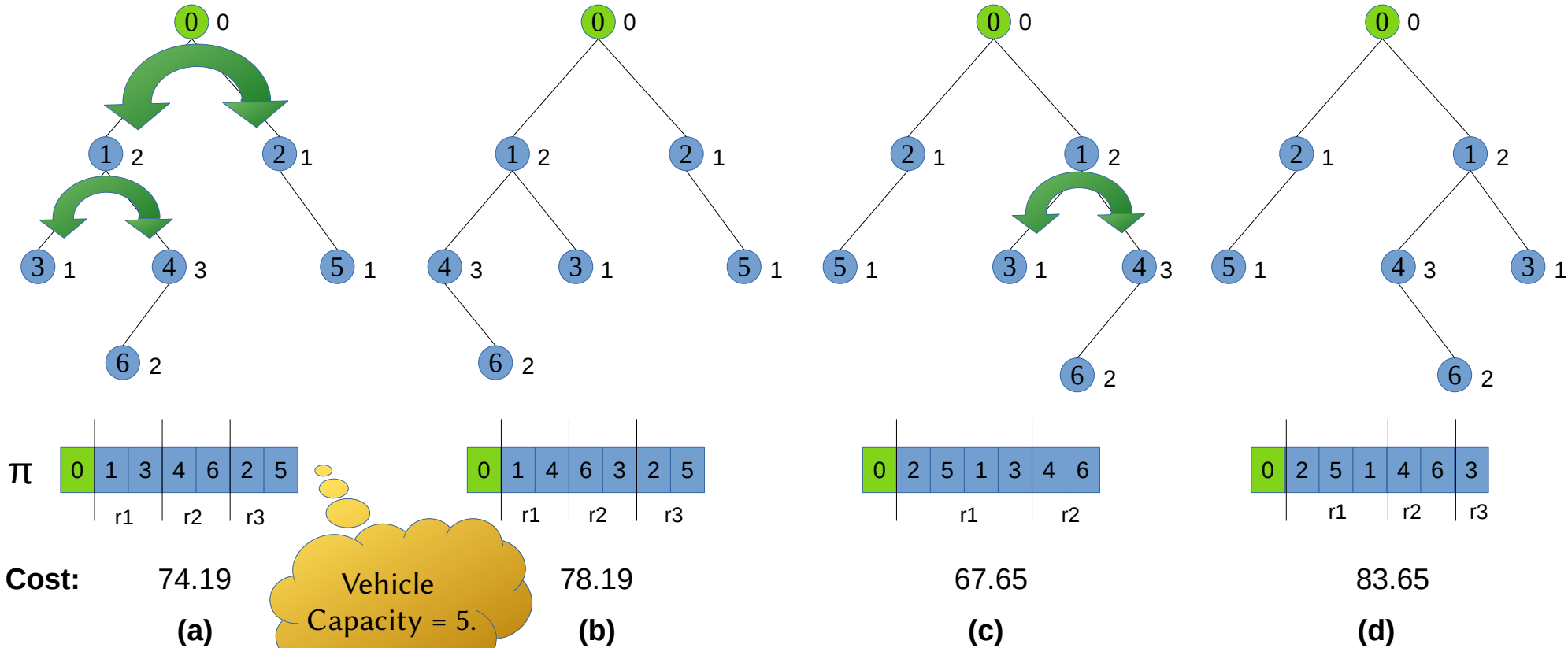
(b) Graph for  $I$ , along with node-demands



(c) Final routes generated by ParMDS

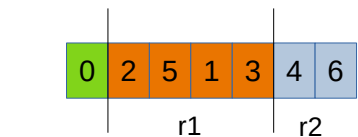
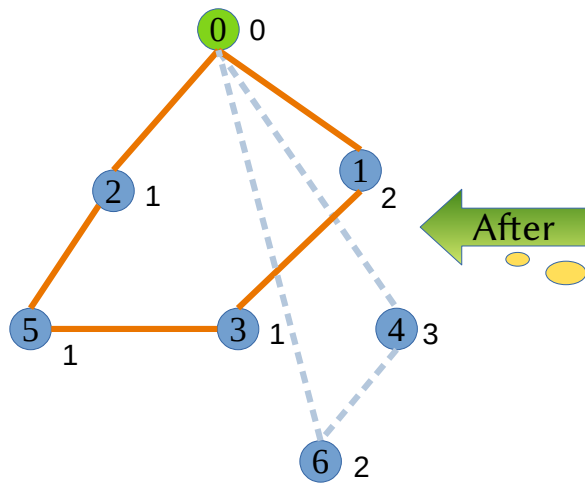
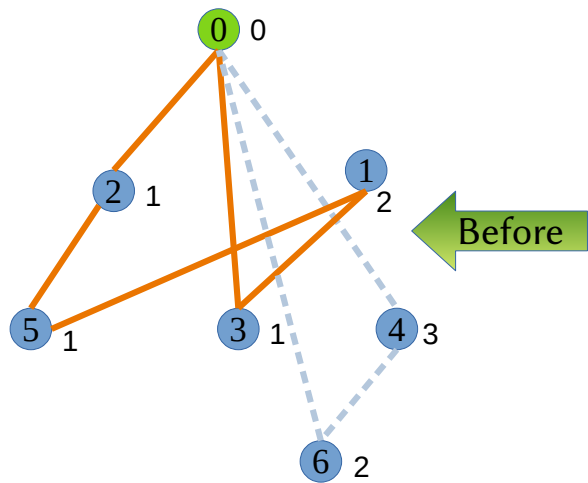
ParMDS on an example input instance with  $n = 7$  and Vehicle Capacity = 5.

# Example - DFS and Randomization



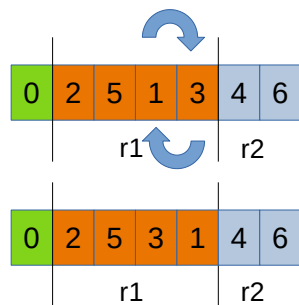
**Takeaway:** Randomizing neighbours of MST may yield a different DFS ordering. Hence, a different route!

# Intra-route optimization - 2Opt



Cost: 67.65

(a)



67.65

↓  
66.13

(b)



# Experiments



- 130 Instances of CVRPLIB

- X
- Golden
- Belgium
- Others

- Intel Xeon CPU E5-2640 v4

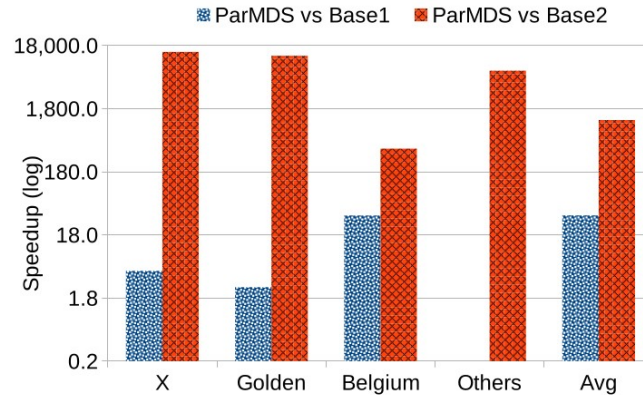
- 40 cores
- Clock 2.4 Ghz
- RAM 64 GB

- Baselines on GPU

- NVIDIA's Tesla P100
- 3584 cores & 12GB global memory
- CUDA 11.5

- Our Code uses

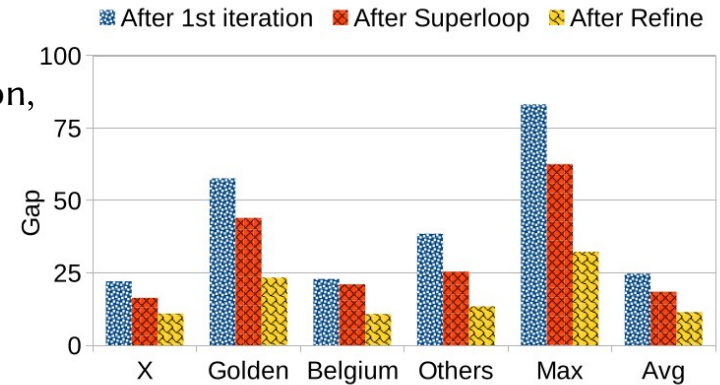
- **SeqMDS:** GCC 9.3.1
- **ParMDS:** nvc++ compiler NVIDIA's HPC SDK 22.11



Method	Execution Time (s) using Random
SeqMDS	1,722.44
ParMDS-Standard	1,522.26
ParMDS-Strided	186.50



Gap at the end of : 1<sup>st</sup> iteration, Superloop and Refine step



# Summary



- GPU parallelization has limitations on larger instances
  - Takes longer time
  - Solution Gap is large
- Our technique combines simpler algorithms/techniques
  - MST and DFS
  - Uses randomization
  - Uses parallelization
  - Open source code
- Our parallelization technique can be extended to other iterative local-search / genetic algorithms



<https://github.com/mrprajesh/parMDS>

## Future Directions

ParMDS can be extended

- to use OpenACC for running on GPUs
- to incorporate direction-aware local-search
- to integrate inter-route optimizations

Thank you

Looking for Postdoc